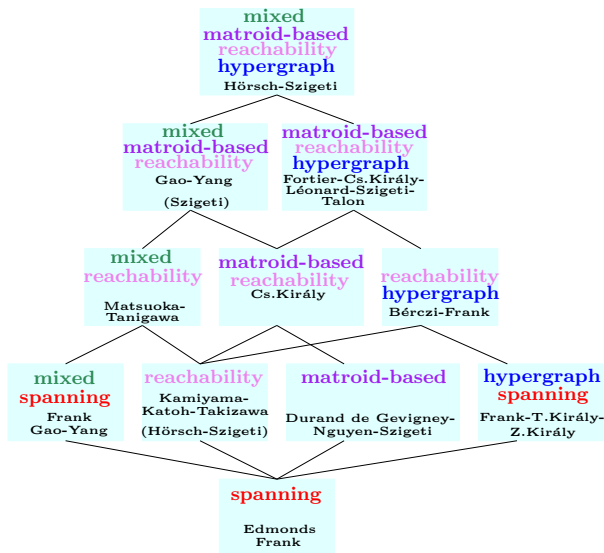


Packing mixed hyperarborescences

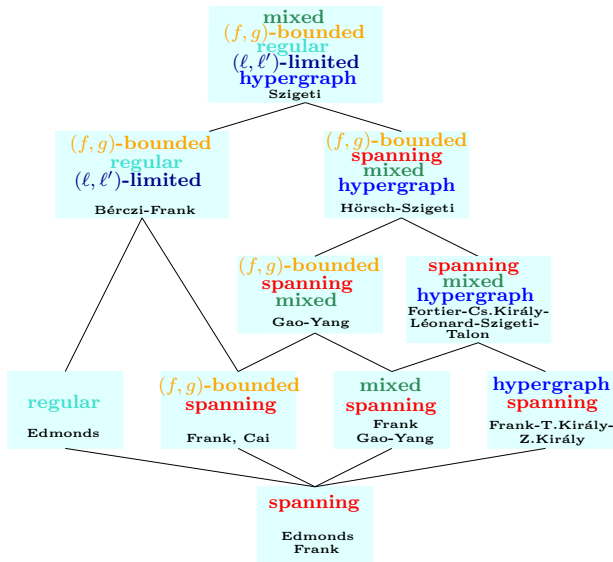
Zoltán Szigeti

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France

Results 1



Results 2



Packing spanning arborescences : fixed roots

Theorem 1 (Edmonds 1973)

Let $D = (V, A)$ be a digraph and S a multiset of V .

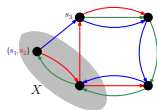
There exists a **packing of spanning s -arborescences** ($s \in S$)

$|S_X| + d_A^-(X) \geq |S|$ for all $\emptyset \neq X \subseteq V$.



Definitions

- 1 **s -arborescence**: unique (s, v) -dipath for all $v \in V$.
- 2 **spanning** subgraph: contains all the vertices.
- 3 **packing** subgraphs: pairwise arc-disjoint subgraphs.
- 4 **S -branching**: unique (S, v) -dipath for all $v \in V$.
- 5 **multiset S** of V : vertex set with multiplicities.
- 6 **S_X** for $X \subseteq V$: restriction of S in X .



Regular packing of arborescences

Definition

k -regular packing of arborescences:
each vertex belongs to k arborescences in the packing.



Theorem 3 (Edmonds 1973)

Let $D = (V, A)$ be a digraph, S a multiset of V and $k \in \mathbb{Z}_+$.

There exists a **k -regular packing of s -arborescences** ($s \in S' \subseteq S$) \iff

$|S_X| + |d_A^-(X)| \geq k$ for all $\emptyset \neq X \subseteq V$.

Remark

There exists a **k -regular** packing of s -arborescences ($s \in S$) \iff

there exists a packing of **k S_i -branchings** with $\bigcup S_i = S$.

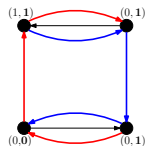
Theorem 3 implies Theorem 1.



(f, g) -bounded packing of spanning arborescences

Definition

(f, g) -bounded packing of arborescences:
number of v -arborescences in the packing is
at least $f(v)$ and at most $g(v) \forall v \in V$.



Theorem 4 (Frank 1978, Cai 1983)

Let $D = (V, A)$ be a digraph, $f, g : V \rightarrow \mathbb{Z}_+$ functions and $k \in \mathbb{Z}_+$.
There exists an (f, g) -bounded packing of k spanning arborescences \iff

- 1 $g(v) \geq f(v)$ for every $v \in V$,
- 2 $\min\{k - f(\overline{UP}), g(UP)\} \geq \sum_{X \in \mathcal{P}} (k - d_A^-(X))$
for every subpartition \mathcal{P} of V .

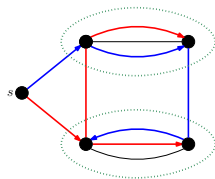
Theorem 4 implies Theorem 2.



Packing spanning mixed arborescences : fixed roots

Definition

- 1 **mixed s -arborescence**:
it can be oriented to obtain an s -arborescence.
- 2 $e_E(\mathcal{P})$: number of edges entering at least one member of a subpartition \mathcal{P} of V .



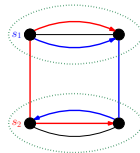
Theorem 5 (Frank 1978)

Let $F = (V, E \cup A)$ be a mixed graph and S a multiset of V .

There exists a **packing of spanning mixed s -arborescences** ($s \in S$) \iff

$e_E(\mathcal{P}) \geq \sum_{X \in \mathcal{P}} (|S| - |S_X| - d_A^-(X))$ for every subpartition \mathcal{P} of V .

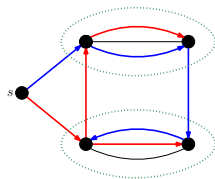
Theorem 5 implies Theorem 1.



Packing spanning mixed arborescences : fixed roots

Definition

- 1 mixed s -arborescence:
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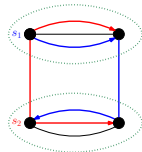
Theorem 5 (Frank 1978)

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Theorem 5 implies Theorem 1.

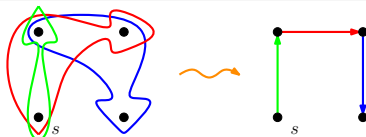


Packing spanning hyperarborescences

Definition

(spanning) (s-)hyperarborescence:

dypergraph that can be trimmed to a (spanning) (s-)arborescence.



Theorem 6 (Frank, T. Király, Z. Király 2003)

Let $\mathcal{D} = (V, \mathcal{A})$ be a dypergraph and S a multiset of V .

There exists a **packing of spanning s-hyperarborescence** ($s \in S$)

$|S_X| + d_{\mathcal{A}}^-(X) \geq |S|$ for all $\emptyset \neq X \subseteq V$.



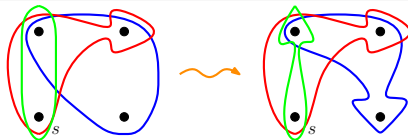
Theorem 6 implies Theorem 1.



Packing spanning mixed hyperarborescences

Definition

(spanning) mixed (s-)hyperarborescence: mixed hypergraph that can be oriented to a (spanning) (s-)hyperarborescence.



Theorem 7 (Fortier, Cs. Király, Léonard, Szigeti, Talon 2018)

Let $\mathcal{F} = (V, \mathcal{E} \cup \mathcal{A})$ be a mixed hypergraph, S a multiset of V .

There exists a **packing of spanning mixed s-hyperarborescence** ($s \in S$) \iff
 $e_{\mathcal{E}}(\mathcal{P}) \geq \sum_{X \in \mathcal{P}} (|S| - |S_X| - d_{\mathcal{A}}^-(X))$ for every subpartition \mathcal{P} of V .

Theorem 7 implies Theorems 5 and 6.

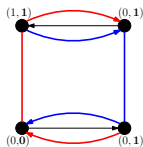


(f, g) -bounded packing spanning mixed arborescences

Theorem 8 (Gao, Yang 2021)

Let $F = (V, E \cup A)$ be a mixed graph, $f, g : V \rightarrow \mathbb{Z}$ functions and $k \in \mathbb{Z}_+$. An (f, g) -bounded packing of k spanning mixed arborescences exists \iff

- 1 $g(v) \geq f(v)$ for every $v \in V$,
- 2 $e_E(\mathcal{P}) + \min\{k - f(\overline{UP}), g(UP)\} \geq \sum_{X \in \mathcal{P}} (k - d_A^-(X))$
for every subpartition \mathcal{P} of V .



Theorem 8 implies Theorems 4 and 5.

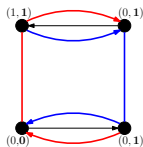


(f, g) -bounded packing spanning mixed hyperarborescences

Theorem 9 (Hörsch-Szigeti 2021)

Let $\mathcal{F} = (V, \mathcal{E} \cup \mathcal{A})$ be a mixed hypergraph, $f, g : V \rightarrow \mathbb{Z}$ functions, $k \in \mathbb{Z}_+$.
An (f, g) -bounded packing of k spanning mixed hyperarborescences exists \iff

- ① $g(v) \geq f(v)$ for every $v \in V$,
- ② $e_{\mathcal{E}}(\mathcal{P}) + \min\{k - f(\overline{\cup \mathcal{P}}), g(\cup \mathcal{P})\} \geq \sum_{X \in \mathcal{P}} (k - d_{\mathcal{A}}^-(X))$
for every subpartition \mathcal{P} of V .



Theorem 9 implies Theorems 8 and 7.



(f, g) -bounded regular (ℓ, ℓ') -limited packing arborescences

Definition

(ℓ, ℓ') -limited packing of arborescences:

packing of at least ℓ and at most ℓ' arborescences.

Theorem 10 (Bérczi, Frank 2018)

Let $D = (V, A)$ be a digraph, $f, g : V \rightarrow \mathbb{Z}_+$ functions and $k, \ell, \ell' \in \mathbb{Z}_+$.
An (f, g) -bounded k -regular (ℓ, ℓ') -limited packing of arborescences exists \iff

- 1 $g_k(v) := \min\{g(v), k\} \geq f(v)$ for every $v \in V$,
- 2 $\min\{g_k(V), \ell'\} \geq \ell$,
- 3 $\min\{\ell' - f(\overline{UP}), g_k(UP)\} \geq \sum_{X \in \mathcal{P}} (k - d_A^-(X))$
for every subpartition \mathcal{P} of V .

Theorem 10 implies Theorems 3 and 4.



Theorem 11 (Szigeti 2022+)

Let $\mathcal{F} = (V, \mathcal{E} \cup \mathcal{A})$ be a mixed hypergraph, $f, g : V \rightarrow \mathbb{Z}$ functions, $k, \ell, \ell' \in \mathbb{Z}_+$.
An (f, g) -bounded k -regular (ℓ, ℓ') -limited packing of mixed hyperarborescences exists \iff

- ① $g_k(v) \geq f(v)$ for every $v \in V$,
- ② $\min\{g_k(V), \ell'\} \geq \ell$,
- ③ $e_{\mathcal{E}}(\mathcal{P}) + \min\{\ell' - f(\overline{\cup \mathcal{P}}), g_k(\cup \mathcal{P})\} \geq \sum_{X \in \mathcal{P}} (k - d_{\mathcal{A}}^-(X))$
for every subpartition \mathcal{P} of V .

Theorem 11 implies Theorems 9 and 10.



Sketch of the proof

- 1 characteristic vectors of the dyperedge sets of the (f, g) -bounded k -regular (ℓ, ℓ') -limited packings of hyperarborescences in orientations of $\mathcal{F} =$ integer points of the intersection of the two generalized polymatroids
 - $\sum_{v \in V} (Q(0, r_v) \cap K(k - g_k(v), k - f(v)))$ and
 - $Q(0, r_{M_{\mathcal{F}}^k}) \cap K(k|V| - \ell', k|V| - \ell)$.
- 2 Frank's theorem on intersection of two generalized polymatroids provides the result.



Open problems

- ① Packing of mixed branchings with given root set sizes ℓ_i .
 - ① For directed graphs: (Bérczi, Frank).
 - ② For undirected graphs: (easy exercise).
 - ③ For each $\ell_i = \ell$: (Szigeti).
- ② Packing of k mixed hyperbranchings each of root set size ℓ .
 - ① For mixed graphs: (previous observation).
 - ② For dypergraphs: (Bérczi, Frank).
 - ③ For hypergraphs: (Martin, Szigeti).

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Thanks for your attention!