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On a min-max theorem on bipartite graphs

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Abstract

Frank et al. (Math. Programming Stud. 22 (1984) 99–112) proved that for any connected bipartite graft (G, T), the minimum size of a T-join is equal to the maximum value of a partition of A, where A is one of the two colour classes of G. Their proof consists of constructing a partition of A of value |F|, by using a minimum T-join F. That proof depends heavily on the properties of distances in graphs with conservative weightings. We follow the dual approach, that is starting from a partition of A of maximum value k, we construct a T-join of size k. Our proof relies only on Tutte's theorem on perfect matchings. It is known (J. Combin. Theory Ser. B 61 (2) (1994) 263–271) that the results of Lovász on 2-packing of T-cuts, of Seymour on packing of T-cuts in bipartite graphs and in grafts that cannot be T-contracted onto $(K_4, V(K_4))$, and of Sebő on packing of T-borders are implied by this theorem of Frank et al. The main contribution of the present paper is that all of these results can be derived from Tutte's theorem.

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1. Introduction

This paper concerns matchings and T-joins. Since T-joins are generalizations of matching, the minimum weight T-join problem contains the minimum weight perfect matching problem. On the other hand, Edmonds and Johnson [2] showed that the former problem can be reduced to the latter one. Thus, these problems are—in fact—equivalent.

In matching theory lots of min-max results are known. Concerning matchings, in fact, we shall consider Tutte's theorem [11] on the existence of perfect matchings in general graphs, and not the min-max version, the Tutte-Berge formula. Concerning

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T-joins, we mention the following min–max theorems: The results of Edmonds–Johnson [2], Lovász [7] on 2-packing of T-cuts, of Seymour [9,10] on packing of T-cuts in bipartite graphs and in grafts that cannot be T-contracted onto $(K_4, V(K_4))$, of Sebő [8] on packing of T-borders and a generalization of Seymour's theorem due to Frank et al. [4]. (For the definitions and the theorems see [3] or [5].) There are some easy known implications between these results, some others can be found in [5], where we showed that the result of Frank et al. [4] implies all of these results, including the Tutte theorem.

Our aim in this paper is to demonstrate a new (surprising) implication, namely, Tutte's theorem implies the result of Frank et al. [4], and consequently, all of these min-max results can be derived from Tutte's theorem.

2. Definitions, notation

In this paper H=(V,E) denotes a graph where V is the set of vertices and E is the set of edges. G=(A,B;E) denotes always a bipartite connected graph and $T\subseteq A\cup B$ a subset of vertices of even cardinality. The pair (G,T) is called a bipartite graft. An edge set $F\subseteq E$ is a T-join if $T=\{v\in A\cup B\colon d_F(v) \text{ is odd}\}$. The minimum size of a T-join is denoted by $\tau(G,T)$. We mention that a bipartite graft (G,T) contains always a T-join.

For a bipartite graft (G = (A, B; E), T) let us introduce an auxiliary graph $G_A := (T, E_A)$ on the vertex set T, where for $u, v \in T$, $uv \in E_A$ if at least one of u and v belongs to A and there exists a path in G connecting u and v of length one or two.

Let K be a vertex set in G. Then $\delta(K)$ denotes the set of edges connecting K and $(A \cup B) - K$. G[K] denotes the subgraph induced by K. b_K^T is defined to be 0 or 1 depending on the parity of $|T \cap K|$. K is called T-odd if $b_K^T = 1$ and T-even if $b_K^T = 0$. For a subgraph K of G, $\bar{K} = G[V(G) - V(K)]$.

We shall need the following operation applied for grafts. For a connected subgraph K of G, by T-contracting K we mean the graft (G',T') obtained from (G,T) where G'=G/K (that is K is contracted into one vertex v_K) and T'=T-V(K) if $b_K^T=0$ and $T'=T-V(K)+\{v_K\}$ if $b_K^T=1$.

In what follows a *component* of a graph means a connected component. For $X \subseteq V(G)$, $\mathcal{K}(G-X)$ denotes the set of components of G-X and $\mathcal{K}_T(G-X)$ denotes the set of T-odd components of G-X. Let $q_T(G-X) = |\mathcal{K}_T(G-X)|$.

We denote by $\mathscr{P}_A := \{u: u \in A\}$ the partition of A where the elements of \mathscr{P}_A are the vertices in A as singletons. The value of a (sub)partition $\mathscr{P} = \{A_1, \ldots, A_k\}$ of A is defined to be

$$val(\mathscr{P}) = \sum \{q_T(G - A_i): A_i \in \mathscr{P}\},\tag{1}$$

in other words,

$$\operatorname{val}(\mathscr{P}) = \sum \left\{ b_K^T : K \in \bigcup_{A_i \in \mathscr{P}} \mathscr{K}(G - A_i) \right\}. \tag{2}$$

The theorem of Frank et al. [4] that generalizes all the min-max results mentioned in the Introduction is as follows.

Theorem 1 (Frank, Sebő, Tardos). If (G,T) is a bipartite graft with G=(A,B;E), then

$$\tau(G, T) = \max\{\text{val}(\mathcal{P}): \mathcal{P} \text{ is a partition of } A\}. \tag{3}$$

In order to be able to prove Theorem 1 by induction we will have to prove a slightly stronger result than Theorem 1. To present it we need some definitions. An edge set C of a connected graph G is called *bicut* if G-C has exactly two connected components. Note that each edge of a tree is a bicut. Let $\mathscr{P}=\{A_1,\ldots,A_k\}$ be a partition of A and let $\mathscr{Q}=\{B_1,\ldots,B_l\}$ be a partition of B. Then $\mathscr{P}\cup\mathscr{Q}$ is called a *bi-partition* of $A\cup B$ in G. Let us denote by $G/(\mathscr{P}\cup\mathscr{Q})$ the bipartite graph obtained from G by identifying the vertices in G for every member G0 and by taking the underlying simple graph. A bi-partition G1 of G2 of G3 is called admissible if

- (i) $F := G/(\mathcal{P} \cup \mathcal{Q})$ is a tree, and
- (ii) for each edge e of F, the edge set of G that corresponds to e forms a bicut of G.

By Claim 4, for any bipartite graft there exists an admissible bi-partition.

Theorem 2. If (G,T) is a bipartite graft with G = (A,B;E), then

$$\tau(G,T) = \max\{\text{val}(\mathcal{P}): \mathcal{P} \cup \mathcal{Q} \text{ is an admissible bi-partition of } A \cup B\}. \tag{4}$$

The proof of Frank et al. [4] for Theorem 1 consists of constructing a partition of A of value |F|, by using a minimum T-join F. That proof depends heavily on the properties of distances in graphs with conservative weightings. We follow the dual approach, that is starting from a bi-partition of $A \cup B$ of maximum value k, we construct a T-join of size k. Our proof applies induction. Taking a special optimal admissible bi-partition either we can use induction for some contracted graphs (and here we need admissibility of the bi-partition) or we can apply Tutte's theorem on perfect matchings, namely a graph H has a perfect matching if and only if $q_V(H-X) \leq |X|$ for every vertex set X of V(H).

We must mention two papers on this topic. Kostochka [6] and Ageev and Kostochka [1] proved results similar to Theorem 2. Their proof technique is different from the present one.

3. Preliminary results

Claim 3. Let (G = (A, B; E), T) be a bipartite graft.

(a) Then the bi-partition $\mathcal{P} \cup \mathcal{Q}$ of $A \cup B$ satisfies (i) where $\mathcal{P} := \{a: a \in A\}$ and $\mathcal{Q} := \{B\}.$

(b) If $X \subseteq A$, then the bi-partition $\mathcal{P} \cup \mathcal{Q}$ of $A \cup B$ satisfies (i) where $\mathcal{P} := \{a: a \in A - X\} \cup \{X\}$ and $\mathcal{Q} := \{K \cap B: K \in \mathcal{K}(G - X)\}.$

The following claim (whose proof is left for the reader) shows that for any bipartite graft there exists an admissible bi-partition.

Claim 4. Let (G = (A, B; E), T) be a bipartite graft.

- (a) If there is no cut vertex in A then $\mathcal{P} \cup \mathcal{Q}$ is an admissible bi-partition of $A \cup B$, where $\mathcal{P} := \{a: a \in A\}$ and $\mathcal{Q} := \{B\}$.
- (b) If there is a cut vertex $v \in A$, that is G can be decomposed into two connected bipartite subgraphs $G_1 = (A_1, B_1; E_1)$ and $G_2 = (A_2, B_2; E_2)$ with exactly one vertex in common, namely v, then let us denote by (G_1, T_1) and (G_2, T_2) the two grafts obtained from (G, T) by T-contracting $V(G_2)$ and $V(G_1)$. If for $i = 1, 2, \mathscr{P}_i \cup \mathscr{Q}_i$ is an admissible bi-partition of $A_i \cup B_i$ and $v \in A'_i$ then $\mathscr{P} \cup \mathscr{Q}$ is an admissible bi-partition of $A \cup B$, where $\mathscr{P} := (\mathscr{P}_1 A'_1) \cup (\mathscr{P}_2 A'_2) \cup \{A'_1 \cup A'_2\}$ and $\mathscr{Q} := \mathscr{Q}_1 \cup \mathscr{Q}_2$.

The definition of an admissible bi-partition implies at once the following claim.

Claim 5. Let $\mathcal{P} \cup \mathcal{Q}$ be an admissible bi-partition of $A \cup B$.

- (a) $K \in \mathcal{K}_T(G A_i)$ for some $A_i \in \mathcal{P}$ if and only if $\bar{K} \in \mathcal{K}_T(G B_i)$ for some $B_i \in \mathcal{D}$.
- (b) $val(\mathcal{P}) = val(\mathcal{Q})$.

Claim 6. Let \mathscr{P} be a partition of A and F a T-join in a bipartite graft (G = (A, B; E), T).

- (a) Then $val(\mathcal{P}) \leq |F|$.
- (b) Moreover, if $val(\mathscr{P}) = |F|$, then for every component K of $G A_i$ for any $A_i \in \mathscr{P}$, $|\delta(K) \cap F| = b_K^T$.

Proof. Let $\mathscr{R} := \bigcup_{A_i \in \mathscr{P}} \mathscr{K}(G - A_i)$. By parity, for each $K \in \mathscr{R}$, $b_K^T \leq |\delta(K) \cap F|$.

Since for $K_1, K_2 \in \mathcal{R}, \delta(K_1) \cap \delta(K_2) = \emptyset$, we have

$$\operatorname{val}(\mathscr{P}) = \sum_{K \in \mathscr{R}} b_K^T \leqslant \sum_{K \in \mathscr{R}} |\delta(K) \cap F| \leqslant |F|. \qquad \Box$$

Claim 7. For every partition \mathscr{P} of A in a bipartite graft (G = (A, B; E), T), $val(\mathscr{P}) \equiv |T \cap A| \pmod{2}$.

Proof. Since
$$|T|$$
 is even, for each $A_i \in \mathscr{P}, q_T(G - A_i) \equiv |T \cap A_i| \pmod{2}$. Thus $\operatorname{val}(\mathscr{P}) = \sum_{A_i \in \mathscr{P}} q_T(G - A_i) \equiv \sum_{A_i \in \mathscr{P}} |T \cap A_i| = |T \cap A|$.

We shall deal with some bi-partitions along the proofs. The admissibility of these bi-partitions can always be easily verified. The following easy fact may be useful.

Claim 8. Let X be a subset of vertices of a connected graph H. Let K be a component of H-X. If X is contained in one of the components of H-K, then H-K is connected.

Claim 9. Let H be a connected graph with |V(H)| even. If X is a minimal vertex set with $q_V(H-X) > |X|$, then for every component K of H-X, H-K is connected.

Proof. By assumption, using the usual parity argument, $q_V(H-X) \ge |X|+2$. Let K be any component of H-X. Then at least one component N of H-K contains more odd components of H-X than vertices in X, that is $q_V(H-(N\cap X)) > |N\cap X|$. Then, by the minimality of X, $N\cap X=X$, that is, by Claim 8, H-K is connected. \square

Claim 10. Let (G = (A, B; E), T) be a bipartite graft. If the auxiliary graph G_A has a perfect matching M then G contains a T-join of cardinality $|T \cap A|$.

Proof. For every edge $uv \in M$ there exists a (u,v)-path in G of length at most two. Since M is a matching these paths are edge disjoint. The union F of these paths is a T-join of G because M covers all the vertices of T. By construction, $|F| = |T \cap A|$. \square

4. The proof of Theorem 2

Let (G,T) be a counterexample with minimum number of vertices in G. By Claim 6(a), for any admissible bi-partition $\mathscr{P} \cup \mathscr{Q}$ of $A \cup B$, $\operatorname{val}(\mathscr{P}) \leqslant \tau(G,T)$, so $\operatorname{val}(\mathscr{P}) < \tau(G,T)$.

Lemma 11. G is 2-connected.

Proof. Suppose that G contains a cut vertex v, by symmetry we may suppose that $v \in A$. We use the notation of Claim 4. For i = 1, 2, (G_i, T_i) is a bipartite graft and $|A_i \cup B_i| < |A \cup B|$ so there exists an admissible bi-partition $\mathscr{P}_i \cup \mathscr{Q}_i$ of $A_i \cup B_i$ with

$$\tau(G_i, T_i) = \text{val}(\mathcal{P}_i). \tag{5}$$

Clearly,

$$\tau(G,T) = \tau(G_1, T_1) + \tau(G_2, T_2). \tag{6}$$

Let $\mathcal{P} \cup \mathcal{Q}$ be the admissible bi-partition of $A \cup B$ defined in Claim 4(b). Note that

$$val(\mathcal{P}) = val(\mathcal{P}_1) + val(\mathcal{P}_2). \tag{7}$$

Then, by (6), (5) and (7), $\tau(G,T) = \text{val}(\mathcal{P})$ showing that (G,T) is not a counterexample. \square

Let us denote by MAX the maximum value of an admissible bi-partition of $A \cup B$. Observe that MAX $\geqslant |T \cap A|$ and MAX $\geqslant |T \cap B|$. The first comes from the admissible bi-partition $\mathscr{P} = \{v: v \in A\}, \mathscr{Q} = \{B\}$, the other one from $\mathscr{P} = \{A\}, \mathscr{Q} = \{v: v \in B\}$. These bi-partitions are admissible by Claim 4(a).

Case 1: First suppose that $MAX = |T \cap A|$ (or $MAX = |T \cap B|$).

Lemma 12. If the auxiliary graph G_A has no perfect matching then there exists an admissible bi-partition $\mathcal{P} \cup \mathcal{Q}$ of $A \cup B$ with val $(\mathcal{P}) > |T \cap A|$.

Proof. By Tutte's Theorem, there exists a set $X \subset T$ so that $q_T(G_A - X) > |X|$. Let us take a minimal such set.

We claim that $X \cap B = \emptyset$. Suppose that $a \in X \cap B$. Suppose that a is connected to two odd components K_1 and K_2 of $G_A - X$. Then, by the definition of G_A , there is an edge between K_1 and K_2 , that is they cannot be different components of $G_A - X$. Thus a is connected to at most one odd component of $G_A - X$. Hence, $q_T(G_A - (X - a)) \ge q_T(G_A - X) - 1 \ge |X| > |X - a|$, contradicting the minimality of X.

Let us denote by B_1 the set of vertices in B-T that has at least one neighbour in $A\cap T$ and let $B_2:=B-T-B_1$. Let $G_1:=G[T\cup B_1]$ and $G_2:=G[(A-T)\cup B_2]$. Note that by the definition of G_A there is a bijection between the components of G_A-X and the components of G_1-X different from isolated vertices in B_1 . Moreover, the T parity of the corresponding components are the same. Let $\mathscr{R}=\mathscr{K}(G_2)$. Note that if $R\in\mathscr{R}$ then there is no edge between $R\cap B_2$ and $A\cap T$. We distinguish two cases.

Case I: First, suppose that $X = \emptyset$ that is $q_T(G_1) \ge 1$, in other words $q_T(G - (A - T)) \ge 1$. Let $\mathcal{R}_1 \subseteq \mathcal{R}$ be a minimal subset of \mathcal{R} so that $q_T(G - A') \ge 1$, where $A' := \bigcup \{R \cap A : R \in \mathcal{R}_1\}$. Let $\mathcal{P} = \{u : u \in A - A'\} \cup \{A'\}$ and let $\mathcal{Q} = \{R \cap B : R \in \mathcal{K}(G - A')\}$. By Claim 3(b), $\mathcal{P} \cup \mathcal{Q}$ satisfies (i). Since $A' \subseteq A - T$, $|(V(G) - A') \cap T|$ is even so $q_T(G - A') \ge 2$ and, by the minimality of \mathcal{R}_1 , each such component has at least one neighbour in every $R \in \mathcal{R}_1$. Since G is 2-connected and for every $R \in \mathcal{R}_1$, G[R] is connected, it follows that for every $D \in \mathcal{K}(G - A')$, G - D is connected, that is (ii) is also satisfied, so $\mathcal{P} \cup \mathcal{Q}$ is an admissible bi-partition and

$$\operatorname{val}(\mathscr{P}) = \sum_{A_i \in \mathscr{P}} q_T(G - A_i) \geqslant \sum_{t \in A - A'} b_t^T + q_T(G - A') \geqslant |T \cap A| + 2.$$

Case II: Secondly, suppose that $X \neq \emptyset$. By the minimality of $X, X \subset V(G')$ where $G' \in \mathcal{K}(G_1)$. Let $\mathcal{R}_1 \subseteq \mathcal{R}$ be a minimal subset of \mathcal{R} so that all the components of G' - X rest in different components of G - A'' - X, where $A'' := \bigcup \{R \cap A: R \in \mathcal{R}_1\}$. Let $\mathcal{P} := \{X \cup A''\} \cup \{u: u \in A - (X \cup A'')\}$ and let $\mathcal{Q} = \{R \cap B: R \in \mathcal{K}(G - X - A'')\}$. By Claim 3(b), $\mathcal{P} \cup \mathcal{Q}$ satisfies (i). For each $R \in \mathcal{R}_1$, G[R] is connected and, by the minimality of \mathcal{R}_1 , R has neighbours in at least two different components of G - X - A''. Moreover, by Claim 9, for each $K \in \mathcal{K}(G' - X)$, G' - K is connected, hence $(G - \bigcup \{R: R \in \mathcal{R}_1\}) - K'$ is connected, where $K' \in \mathcal{K}(G - X - A'')$ that contains K. It follows that $X \cup A''$ is contained in one of the components of G - K'. Thus, by Claim 8 and by 2-connectivity,

 $\mathcal{P} \cup \mathcal{Q}$ is an admissible bi-partition of $A \cup B$ and

$$\operatorname{val}(\mathscr{P}) = \sum_{A_i \in \mathscr{P}} q_T(G - A_i) = \sum_{t \in A - X - A''} b_t^T + q_T(G - (X \cup A''))$$
$$= |A \cap T| - |X| + q_T(G_A - X) > |T \cap A|. \quad \Box$$

By Lemma 12, G_A (G_B , resp.) has a perfect matching and thus, by Claim 10, G contains a T-join of cardinality $|T \cap A|$ ($|T \cap B|$, resp.). By Claim 6(a), the proof of the theorem is complete.

Case 2: Secondly, suppose that MAX > $|T \cap A|$ and MAX > $|T \cap B|$. Then, by Lemma 11, every optimal admissible bi-partition contains a set A_i with $1 < |A_i| < |A|$. Let us choose an optimal admissible bi-partition $\mathscr{P} \cup \mathscr{Q}$ of $A \cup B$ so that such a set A_i of \mathscr{P} is as large as possible. Let $K \in \mathscr{K}(G - A_i)$ so that $|V(K)| \geqslant 2$. (Since $|A_i| < |A|$ such a set exists.) Then, by Claim 5, $\overline{K} \in \mathscr{K}(G - B_j)$ for some $B_j \in \mathscr{Q}$ and $|V(\overline{K})| \geqslant 2$. Let us denote by (G_1, T_1) and (G_2, T_2) the two bipartite grafts obtained from (G, T) by T-contracting the connected subgraphs K and \overline{K} , respectively. The colour classes of G_r will be denoted by A^r and B^r , while the contracted vertex of G_r is denoted by V_r for r = 1, 2. Let $\mathscr{P}_1 := \{A_k \in \mathscr{P} : A_k \subseteq A^1\}$ and $\mathscr{Q}_1 := \{B_l \in \mathscr{Q} : B_l \subseteq B^1\} \cup \{v_1\}$. Let $\mathscr{P}_2 := \{A_k \in \mathscr{P} : A_k \subseteq A^2\} \cup \{v_2\}$ and $\mathscr{Q}_2 := \{B_l \in \mathscr{Q} : B_l \subseteq B^2\}$. The admissibility of the bi-partition $\mathscr{P} \cup \mathscr{Q}$ implies the following Claim.

Claim 13.

- (a) $\mathscr{P}_r \cup \mathscr{Q}_r$ is an admissible bi-partition of $A^r \cup B^r$ in $G_r, r = 1, 2$.
- (b) $\operatorname{val}_{(G,T)}(\mathscr{P}) = \operatorname{val}_{(G_1,T_1)}(\mathscr{P}_1) b_{v_1}^{T_1} + \operatorname{val}_{(G_2,T_2)}(\mathscr{P}_2).$

Lemma 14. For $r = 1, 2, \mathcal{P}_r \cup \mathcal{Q}_r$ is an optimal admissible bi-partition of $A^r \cup B^r$ in (G_r, T_r) .

Proof. By Claim 13(a), only the optimality must be verified. By symmetry, it is enough to prove it for r=2. Suppose that $\mathscr{P}'\cup\mathscr{Q}'$ is an admissible bi-partition of $A^2\cup B^2$ in G_2 with $\operatorname{val}_{(G_2,T_2)}(\mathscr{P}')>\operatorname{val}_{(G_2,T_2)}(\mathscr{P}_2)$. Let us denote by X that member of \mathscr{P}' that contains v_2 . Since $\mathscr{P}_1\cup\mathscr{Q}_1$ and $\mathscr{P}'\cup\mathscr{Q}'$ are admissible bi-partitions and \bar{K} is connected, $\mathscr{P}'':=(\mathscr{P}_1-A_i)\cup(\mathscr{P}'-X)\cup\{(X-v_2)\cup A_i\},\ \mathscr{Q}''=(\mathscr{Q}_1-\{v_1\})\cup\mathscr{Q}'$ is an admissible bi-partition of $A\cup B$ in G. By Claim 13(b),

$$\begin{aligned} \operatorname{val}_{(G,T)}(\mathscr{P}'') &= \operatorname{val}_{(G_1,T_1)}(\mathscr{P}_1) - b_{v_1}^{T_1} + \operatorname{val}_{(G_2,T_2)}(\mathscr{P}') \\ &> \operatorname{val}_{(G_1,T_1)}(\mathscr{P}_1) - b_{v_1}^{T_1} + \operatorname{val}_{(G_2,T_2)}(\mathscr{P}_2) = \operatorname{val}_{(G,T)}(\mathscr{P}), \end{aligned}$$

a contradiction.

Lemma 15. If K is T-odd, then for every edge v_2u of G_2 , $\mathscr{P}_2 \cup \mathscr{Q}_2$ is an optimal admissible bi-partition of $A^2 \cup B^2$ in (G_2, T_2') of value $\operatorname{val}_{(G_2, T_2)}(\mathscr{P}_2) - 1$, where $T_2' := T_2 \oplus \{v_2, u\}$.

Proof. By Claim 13(a), only the optimality must be verified. $\operatorname{val}_{(G_2,T_2')}(\mathscr{P}_2) = \operatorname{val}_{(G_2,T_2)}(\mathscr{P}_2) - 1$ because for a component L of $G_2 - R$ with $R \in \mathscr{P}_2 - \{v_2\}$, $|L \cap T_2| \equiv |L \cap T_2'|$ (mod 2) and the unique component K of $G_2 - v_2$ becomes T_2' -even. Suppose that $\mathscr{P}' \cup \mathscr{L}'$ is an admissible bi-partition of $A^2 \cup B^2$ in (G_2, T_2') with $\operatorname{val}_{(G_2, T_2')}(\mathscr{P}') > \operatorname{val}_{(G_2, T_2)}(\mathscr{P}_2) - 1$. By Claim 7, $\operatorname{val}_{(G_2, T_2')}(\mathscr{P}') \geqslant \operatorname{val}_{(G_2, T_2)}(\mathscr{P}_2) + 1$. Note that since K is T-odd, $b_{v_1}^{T_1} = 1$. Let us denote by K that member of \mathscr{P}' that contains V_2 . Since K and K are connected, $\mathscr{P}'' := (\mathscr{P}_1 - A_i) \cup (\mathscr{P}' - X) \cup \{(X - v_2) \cup A_i\}, \mathscr{L}'' = (\mathscr{L}_1 - \{v_1\}) \cup \mathscr{L}'$ is an admissible bi-partition of $K \cup K$ in K.

If $X = v_2$ then, by Claim 13(b),

$$\operatorname{val}_{(G,T)}(\mathscr{P}'') = \operatorname{val}_{(G_1,T_1)}(\mathscr{P}_1) + \operatorname{val}_{(G_2,T_2')}(\mathscr{P}')$$

$$\geq \operatorname{val}_{(G_1,T_1)}(\mathscr{P}_1) + \operatorname{val}_{(G_2,T_2)}(\mathscr{P}_2) + 1 > \operatorname{val}_{(G,T)}(\mathscr{P}),$$

a contradiction.

If $X \neq v_2$, then, by Claim 13(b),

$$val_{(G,T)}(\mathscr{P}'') \ge (val_{(G_1,T_1)}(\mathscr{P}_1) - 1) + (val_{(G_2,T_2')}(\mathscr{P}') - 1)$$

$$\ge val_{(G_1,T_1)}(\mathscr{P}_1) - 1 + val_{(G_2,T_2)}(\mathscr{P}_2) = val_{(G,T)}(\mathscr{P}),$$

that is $\mathscr{P}'' \cup \mathscr{Q}''$ is an optimal admissible bi-partition of $A \cup B$ in G, but $|(X - v_2) \cup A_i| > |A_i|$, contradicting the maximality of A_i . \square

By induction ($|V(G_1)| < |V(G)|$ because |V(K)| > 2) and by Lemma 14, there exists a T_1 -join F_1 in G_1 with $|F_1| = \text{val}(\mathcal{P}_1)$.

First suppose that K is a T-even component of $G-A_i$. By induction $(|V(G_2)| < |V(G)|)$ because $|V(\bar{K})| \ge 2$ and by Lemma 14, there exists a T_2 -join F_2 in G_2 with $|F_2| = \operatorname{val}(\mathscr{P}_2)$. Then, by Claim 6(b), $|F_1 \cap \delta(K)| = 0 = |F_2 \cap \delta(K)|$, hence $F := F_1 \cup F_2$ is a T-join and, by Claim 13(b), $|F| = |F_1| + |F_2| = \operatorname{val}(\mathscr{P}_1) + \operatorname{val}(\mathscr{P}_2) = \operatorname{val}(\mathscr{P})$. By Claim 6(a), we are done.

Now suppose that K is a T-odd component of $G - A_i$. Then, by Claim 6(b), $|F_1 \cap \delta(K)| = 1$. This edge corresponds to an edge v_2u in G_2 . By induction $(|V(G_2)| < |V(G)|$ because $|V(\bar{K})| \ge 2$) and by Lemma 15 with edge v_2u , there exists a T_2 -join F_2 in G_2 with $|F_2| = \text{val}(\mathscr{P}_2) - 1$. Then $F := F_1 \cup F_2$ is a T-join and, by Claim 13(b), $|F| = |F_1| + |F_2| = \text{val}(\mathscr{P}_1) + \text{val}(\mathscr{P}_2) - 1 = \text{val}(\mathscr{P})$. By Claim 6(a), we are done. \square

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