Search techniques in Constraint Programming

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Outline

Backtrack Search

- Variable Ordering
- Value Ordering (Branching)
- Restarts & Randomization

2 Discrepancy Search techniques

Search Algorithms

Tree Search

- Given a decision (a constraint) for example x = v, either:
 - * there exists a solution where x = v
 - * there exists a solution where $x \neq v$
 - \star there is no solution
- ► Typical decisions (fix a value, split a domain, enforce a precedence, ...)
- Exploring both branches gives a complete algorithm
 - In practice, it is often possible to detect inconsistencies after assigning only a few variables
 - * Before each decision, there is a propagation step to enforce local consistencies

Variable Ordering

• The order in which variables are explored matters!

- During search, most of the time is spent in <u>unsatisfiable</u> sub-trees
- Detecting failure early (fail first)
- Start with variables that are most likely to "fail"
 - * Smallest domain size (less freedom, smaller branching factor)
 - * Maximum degree (most constrained variable)
 - ★ Minimum domain size degree
 - * ...

Coloring



















































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Weighted Degree Heuristic

- The weight of a constraint is initialised to 1
- When propagating, if a contradiction is detected on that constraint, its weight is increased by 1
- The <u>weight</u> of a variable is equal to the sum of the weights of its neighboring constraints

weight(x) =
$$\sum_{c \ s.t. \ x \in \mathcal{X}(c)} \text{weight}(c)$$

- Domain over weigthed degree:
 - ► The variable that minimizes the ratio domain size weight is selected first

Impact

- The impact I(x, v) of a pair (x, v) is initialised to 0
- The size of the search space of a network P = (X, D, C) is bounded by ∏_{x∈X} |D(x)|
- When the decision x = v is succesful
 - Let b be the size of the search space of P = (X, D, C) (before the decision)
 - Let a be the size of the search space of P' =AC(X, D, C ∪ (x = v)) (after the decision)

 - The impact of (x, v) is the average impact computed for all decisions x = v.
- The weight of a variable x is $\sum_{v \in D(x)} \frac{1}{I(x,v)}$, the variable with minimum weight is selected first

Value Ordering

• Often considered less important than variable ordering

- Most of the search effort is spent in unsatisfiable subtrees
- Value ordering does not really matter in an unsatisfiable subtree since it must be fully explored
- In satisfaction problems:
 - Choosing a satisfiable subtree when branching (promise: choose the least constrained value)
- In optimisation problems:
 - Choosing the assignment that is likely to give the best objective value in order to get good upper bounds quickly

Heavy Tail Behavior

- Given a collection of instances of a problem, we often observe a few exceptionally hard instances
 - ► These instances are rare, but take exceptionally longer to solve
 - Large impact on the mean runtime for a given set
 - As opposed to normal distributions, the mean does not stabilize when the size of the sample grows
 - \star When the sample grows, the mean runtime is skewed up
 - * Heavy tail behavior
- Not a characteristic of the instance, the same behavior behavior is observed if we run several times the same instance while varying some parameter of the solver
 - For some combination instance / solver parameters, we get trapped into an exponential subtree

Heavy Tail Behavior

- Randomization:
 - Add some randomized parameter in variable or value selection (for instance to break ties)
 - Given the same random seed the solver will explore the same tree, however it will never explore two identical subproblems in the same way
- Restarting:
 - ► After a given limit *r*, for instance in number of explored nodes: restart from scratch
- Randomization + restarts <u>eliminates</u> the huge variance in solver performance
 - And therefore reduces the mean runtime when a heavy tail behavior could be observed

- Which limit *r* should we use?
 - Geometric: $r_i = f^i b$

***** b = 100, f = 2: 100, 200, 400, 800, ...

► Luby:

*
$$r_i = 2^{i-1}b$$
 if $i = 2^k - 1$

★
$$r_{i-2^{k-1}+1}b$$
 if $2^{k-1} \le i \le 2^k - 1$

***** b = 10: 10,10,20,10,10,20,40,10,10,20,10,10,20,40,80 ...

Outline



2 Discrepancy Search techniques

Limits of heuristics

- Idea of Limited Discrepancy Search:
 - A heuristic is never perfect
 - But if we trust it, we want to search first <u>as close as possible</u> to the advised decisions.
- Idea of Depth Bounded Discrepency Search:
 - Early choices are less informed
 - A heuristic is generally less reliable at the top of the tree

Limited Discrepancy Search

