## FORCING COLORATIONS AND THE STRONG PERFECT GRAPH CONJECTURE

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**tract:** We give various reformulations of the Strong Perfect Graph Conjecture, based on a study of coloring procedures, uniquely colorable subgraphs and  $\omega - 1$ -cliques in minimal imperfect graphs.

#### troduction

is a graph  $\omega = \omega(G)$  denotes the cardinality of a maximum clique and  $\alpha = \alpha(G)$  s the cardinality of a maximum stable set.  $\chi = \chi(G)$  is the chromatic number of G. ph means induced subgraph in this paper.

s replace  $\{x\}$  by x throughout the paper. Paths an circuits go through every vertex once. They will be considered to be subgraphs or edge-sets. The vertex-set of the will be denoted by V(G), the edge-set by E(G).

oh G is called perfect if  $\chi(H) = \omega(H)$  for every subgraph H of G, otherwise it imperfect. Polyhedra of the form  $P = \{x \geq 0 : Ax \leq 1\}$ , where A is a 0-1 ave integer vertices if and only if the rows of A are the characteristic vectors of and the vertices of P are the characteristic vectors of stable sets of a perfect graph. straightforward consequence of Lovász's Perfect Graph Theorem (1972 a) and a stheorem on antiblocking (1970,1971), as it was observed by Chvátal (1975).) act Graphs provide a pleasant reformulation of a wide range of integer programs stageous properties with respect to Optimization.

Perfect Graph Theorem has become the material of undergraduate textbooks, has decreased toward the structure of perfect or minimal imperfect graphs in Fulkerson's Lovász's and Padberg's work (1970, 1972 a,b, 1974). A graph is mal imperfect if it is not perfect, but all its subgraphs are perfect.

characterization of perfectness immediately follows from Lovász (1972 b), but

coNP charaterization would be given by the following conjecture of Berge (1961), (1962):

A graph isomorphic to an odd circuit of length at least five is called a *hole*, and the complement of such a graph is called an *antihole*.

Berge's Strong Perfect Graph conjecture (SPGC): If a graph is minimal imperfect, then it is a hole, or an antihole.

Since a co-NP characterization exists already, there could be doubts about whether this conjecture is interesting enough. However, its investigation may be justified by the closed relation it could have with the recognition of perfect graphs. It is of course a nice structural property interesting for its own sake.

Since the proof of the Strong Perfect Graph Conjecture occured to be difficult, researchers of the field have turned toward proving it for subclasses of graphs, and developped algorithms for recognizing subclasses of perfect graphs. A relatively small number of papers deals with the Strong Perfect Graph Conjecture in general.

Let us summarize one direction of general results, and on the way we define some notions. In the sequal we shall suppose that the reader is familiar with these notions and statements, and we shall use them without reference.

Lovász's theorem on partitionability(1972b), (1984):

Minimal imperfect graphs have  $\alpha\omega + 1$  vertices;

It follows then by an easy computation, that

deleting an arbitrary vertex in a minimal imperfect graph, the remaining graph can be partitioned into  $\alpha$   $\omega$ -cliques, and also into  $\omega$   $\alpha$ -stable-sets.

Graphs G with  $\alpha\omega + 1$  vertices and the property that deleting an arbitrary vertex the remaining graph can be partitioned into  $\alpha$   $\omega$ -cliques, and also into  $\omega$   $\alpha$ -stables are called partitionable, or  $(\alpha, \omega)$  graphs.

Since partitionable graphs are easily seen to be imperfect, the above theorem of Lovász provides a NP-characterization of imperfectness. It also implies Lovász's Perfect Graph Theorem stating that

A graph is perfect if and only if its complement is perfect.

Padberg's corollaries (1974): Let G be an  $(\alpha, \omega)$  graph. Let us mention four properties proved by Padberg which will be the most important for us:

- a.) For every  $\omega$ -clique K there exists a unique  $\alpha$ -stable set disjoint from it, let us denote it by S(K); for every  $\alpha$ -stable set S, there exists a unique  $\omega$ -clique disjoint from it, denote it by K(S).
- b.) For every  $v \in V(G)$ , G v has one unique partition into  $\omega$ -cliques, and one unique

into  $\alpha$ -stable sets. The former partition is the clique-partition of G - v, its are the clique classes of G - v, whereas the latter partition is the coloration of s elements are the color classes of G - v.

 $v \in V(G)$  is contained in  $\alpha$   $\alpha$ -stable sets, namely

K is a clique-class of G-v }, and  $\omega$   $\omega$ -cliques, namely

S is a color-class of G - v \.

paper the representation of a set with a vector will always mean its characteristic a vertex-set.

set of  $\omega$ -cliques of a minimal imperfect graph is linearly independent.

### 's remark on transversals:

d (1976), (1984) has made a crucial remark in the interior of a proof:

inimal imperfect graph G there is no set of vertices whose cardinality is  $\alpha + \omega - 1$  sects every  $\alpha$ -stable set and every  $\omega$ -clique.

an immediate consequence just of the perfectness of subgraphs with  $(\alpha - 1)(\omega - 1)$ 

whose cardinality is  $\alpha + \omega - 1$  and intersects every  $\alpha$ -stable set and every  $\omega$ -clique alled a *small transversal*.

### erization of unique colorability

h which has one unique partition into  $\omega = \omega(G)$  stable sets, is called uniquely Deleting a vertex from a minimal imperfect graph we get a uniquely colorable raph (Padberg). The significance of uniquely colorable perfect graphs, and their o the Strong Perfect Graph Conjecture was first recognized by A. Tucker (1984). conjectures about the relation of what he called "sequential colorings" and unique by, and pointed out the role such a relation could play in a proof of the Strong raph Conjecture.

proach was further developped in Fonlupt, Sebő (1990), where a good characteruniquely colorable perfect graphs is given, and the following separation of the two parts is exhibited:

- 2. Prove the existence of "combinatorial forcings" in uniquely colorable perfect
- 2. Prove that a minimal imperfect graph with "enough combinatorial forcings" antihole.

1 to all graphs G - x ( $x \in V(G)$ ), where G is the minimal imperfect graph of Problem 2, one gets a sufficient number of forcings to apply the answer to Problem 2.

Thus the more particular the forcings in Problem 1 are, or the weaker the condition of Problem 2 is, the closer we are to the SPGC.

The main goal of this work is to make a step forward in Problem 2.

Let G be an arbitrary graph. Let us say that xKy  $(x \neq y \in V(G), K \subseteq V(G))$  is a forcing (to color x and y to the same color in every  $\omega$ -coloration), if both  $K \cup x$  and  $K \cup y$  are  $\omega$ -cliques. We shall say that x and y are forced, if there exist forcings  $x_1K_1x_2$ ,  $x_2K_2x_3$ , ...,  $x_{t-1}K_{t-1}x_t$ ,  $x = x_1$ ,  $y = x_t$ . Indeed, if  $\chi(G) = \omega(G)$ , and x, y are forced, then clearly, x and y must have the same color in every  $\omega$ -coloration of G.

Forcings of the complement of a graph will be called *co-forcings*, and pairs of vertices forced in the complement will be called *co-forced*.

Chvátal (1976), (1984), Giles, Trotter Tucker (1984), Tucker (1984) can be reformulated as statements proving that minimal imperfect graphs with some forcings are holes and antiholes. These reformulations, and some other statements of the kind appear in Bacsó (1989).

S.E. Markossian, G.S. Gasparian and A.S. Markossian's result (1986) is a breakthrough in this direction:

If in G two adjacent vertices are forced, then obviously,  $\chi(G) > \omega(G)$ , in particular G is not perfect. But does G contain a hole or an antihole?

In a hole, any two adjacent vertices are co-forced, and any two vertices at distance two are forced. Markossian and Gasparian and Markossian (1986) prove that this statement can be reversed:

**Theorem 1.1** If a minimal imperfect graph G and its complement both have two forced vertices which are adjacent, then G is a hole or antihole.

The key-result of the present work is the following generalization:

Theorem 1.2 If a minimal imperfect graph G has two forced vertices which are adjacent, then G is a hole or antihole.

For a proof of the SPGC it suffices now to prove that a minimal imperfect graph or its complement have adjacent forced vertices.

We prove Theorem 1.2 in Section 2. This provides a new, short proof of Theorem 1.1 itself. (In Markossian and Gasparian and Markossian's solution an ordered sequence of  $\alpha$  different forcings, and  $\omega$  different co-forcings is essential, their proof is quite involved, and

uses an exhaustive case checking; the proof in Section 2 makes use of one forcing and two co-forcings only, and exploits these in a simple way.)

The other pole of this paper is a study of the  $\omega-1$ -cliques of a minimal imperfect graph, with the goal of understanding forcings better, and getting closer to the conjectures below. The main results about  $\omega-1$ -cliques are sketched in Section 4. We could still not be a grown the existence of at least one forcing in a minimal imperfect graph, which would be a preakthrough we think, but an interesting structure seems to come up: we shall call some estricted kind of  $\omega-1$ -cliques intervals, and our goal is to prove that intervals arise as the intersection of  $\omega$ -cliques. (The name "interval" comes from the fact that specializing hem to "webs" they are exactly the  $\omega-1$ -cliques forming intervals, and the properties he shall prove also remind intervals.) This turns out to be a rich notion: in Section 4 we shall prove also remind intervals, intervals and  $\omega$ -cliques. Our main tool for this ork out the relations between vertices, intervals and  $\omega$ -cliques. Our main tool for this ork will be Theorem 1.2. We shall also prove and then use a new relation between unique lorability and forcings (that is some results on Problem 1), which may be interesting for own sake:

**neorem 1.3** Let G be partitionable and suppose K is an  $\omega$ -clique, S is an  $\alpha$ -stable,  $TK \cap S = \emptyset$ . If  $G - K \cup S$  and its complement are uniquely colorable, then G is a hole antihole.

proof of this theorem can be found at the end of Section 3.

he results of the paper lead us into a new range of simple reformulations of the Perfect ph Conjecture, which we summarize in Section 5.

t us finish this introduction with two conjectures, which imply together the Strong ect Graph conjecture, and which are the main motivation for the present research. first is a weakening of a conjecture of Tucker (1984), see also Fonlupt, Sebő (1990). second would sharpen the results of the paper.

ecture 1 If G is a uniquely colorable perfect graph such that  $\alpha(G)\omega(G) = V(G)$  be set of its  $\omega$ -cliques is linearly independent, then it has a forcing.

cture 2 If G is a partitionable graph which has both a forcing and a co-forcing, also has a small transversal.

### 2. Critical edges and the Perfect Graph Conjecture

An edge  $e \in E(G)$  is called *critical*, if  $\alpha(G - e) = \alpha + 1$ . This means exactly that there exists S,  $|S| = \alpha - 1$  such that xSy is a co-forcing.

If the pair (x, y) is a critical edge of the complement we shall call it a critical co-edge. (There exists then a unique  $\omega - 1$ -clique K so that xKy is a forcing.)

Throughout this section we suppose that G is an  $(\alpha, \omega)$  partitionable graph.

Markossian, Gasparian and Markossian (1988) made several simple but important observations about critical edges. Let us state three of them:

(2.1) If  $x_0x_1, \ldots, x_{k-1}x_k$  are critical edges,  $k < \omega$ , then  $\{x_0, \ldots, x_k\}$  is a clique.

Indeed, let  $x_{i-1}S_ix_i$  be the coforcings corresponding to these critical edges (i = 1, ..., k). Then  $\{x_0, ..., x_k\} \cup S_1 \cup ... \cup S_k$  is a proper subset of the vertices of G, because  $k+1+k(\alpha-1)=k\alpha+1<\omega\alpha+1$ . But a proper subset of an  $(\alpha,\omega)$ -graph has a partition into at most  $\alpha$   $\omega$ -cliques, and coforced vertices must be in the same clique-class, whence they are adjacent, as claimed.

(2.2) If xSy is a co-forcing, then there exists a unique  $\omega$ -clique  $K_x$  containing x and not containing y, and  $K_x = K(S \cup y)$ .

*Proof.* Clearly, an  $\omega$ -clique  $K_x$  containing x is disjoint from S. If in addition  $y \notin K_x$ , then  $K_x \cap (S \cup y) = \emptyset$ .

Let us first prove that Theorem 1.2 is a special case of the following Theorem 2.1. Indeed, if there are co-forced non-adjacent vertices, then it is easy to see from (2.1) that there exists a path  $(x_0, \ldots, x_{\alpha})$ , where  $x_{i-1}x_i$   $(i=1,\ldots,\alpha)$  are critical edges. Applying (2.1) again, we immediately get that  $x_0Kx_{\alpha}$  is a forcing, where  $K=\{x_1,\ldots,x_{\alpha-1}\}$ . This forcing and the two critical edges  $x_0x_1, x_{\alpha-1}x_{\alpha}$  satisfy the condition of the theorem below.

**THEOREM 2.1** Suppose G is an  $(\alpha, \omega)$  graph, and it has no small transversal. If there exists  $v_1, v_2 \in V(G)$  and an  $\omega$ -clique K such that  $v_1 K v_2$  is a forcing, moreover there exist  $u_1, u_2 \in K$  (not necessarily distinct), such that  $u_1 v_1, u_2 v_2$  are critical edges, then G is a hole or an antihole.

Proof. Let  $u_1S_1v_1$ ,  $u_2S_2v_2$  be the two co-forcings corresponding to the two critical edges of the theorem. If  $u_1 = u_2$ , then by (2.1)  $\omega = 2$ . (Because  $v_1v_2 \notin E(G)$ , but there is a path consisting of critical edges between them.) Thus G is a hole.

Suppose now  $u_1 \neq u_2$  throughout the proof.

Let S be the family of  $\alpha$ -stable sets containing  $u_1$ , different from  $S_1 \cup u_1$ , and not a color class of  $G - u_2$ . We have  $|S| = \alpha - 2$ .

Case 1 There exists  $S \in S$  such that  $S \setminus u_1$  is not a subset of  $S_1 \cup S_2$ .

Let then  $s \in (S \setminus u_1) \setminus (S_1 \cup S_2)$ , and let Q be the clique class of G - s for which  $v_2 \in Q$ . (Equivalently,  $s \in S(Q)$ ,  $v_2 \in Q$ .)

Let  $T := (S \setminus u_1) \cup (Q \setminus v_2) \cup v_1$ .

 $Q \cap S \neq \emptyset$ , because otherwise S would be a color class of  $G - v_2$ .  $|T| = \alpha + \omega - 2$  follows.

We first show that every  $\omega$ -clique R different from K(S) has a non-empty intersection with T. Indeed, since  $R \cap S \neq \emptyset$ , this is obvious if  $u_1 \notin R$ ; if  $u_1 \in R$ , then either  $v_1 \in R$  and we are done, or  $v_1 \notin R$ , and then by (2.2)  $R = K(S_1 \cup v_1) = K \cup v_2$ . But then  $C \in K \cap T$ :  $C \in K$  by assumption;  $C \in K \cap T$ :  $C \in K$  by assumption;  $C \in K \cap T$ :  $C \in K$  by assumption of  $C \in K$ .

Second, we show that every  $\alpha$ -stable set U different from  $S_2 \cup v_2$  has a non-empty versection with T. Indeed,  $s \in S(Q) \cap T$ . If  $U \neq S(Q)$ , then  $(Q \setminus v_2) \cup v_1$  intersects it, tept if Q is the unique stable set (see (2.2)) containing  $v_2$  and not containing  $v_1$ . But a stable set is just  $S_2 \cup v_2$ .

ince  $S \neq S_2 \cup v_2$ , K(S) and  $S_2 \cup v_2$  are not disjoint. Adding their intersection to T we a small transversal, a contradiction. Thus Case 1 cannot hold.

ut then Case 1 can also not hold if we interchange the indices 1 and 2. The only case remains:

**e** 2 Every  $\alpha$ -stable set which contains  $u_1$  or  $u_2$  and is not a color class of  $G - u_1$  or  $u_2$  is a subset of  $S_1 \cup S_2 \cup \{u_1, u_2\}$ .

H be the graph induced by  $S_1 \cup S_2 \cup \{u_1, u_2\}$ . By the assumption,  $S_1 \cup S_2 \cup \{u_1, u_2\}$  ins all the  $\alpha$ -stable sets containing  $u_1$  except one, and the same is true for  $u_2$ . Thus H has at least  $2(\alpha - 1)$   $\alpha$ -stable sets.

show that this is possible only if  $\alpha = 2$ . Let H be the graph induced by  $S_1 \cup S_2$ . Since ipartite, it is perfect, whence it can be partitioned to  $\alpha$  cliques, that is, there exists a matching  $M = \{x_1y_1, \ldots, x_{\alpha}y_{\alpha}\}$  in H. (We implicitly used here the Perfect Graph m, or König's theorem about matching's in bipartite graphs. To use the latter note  $H = \alpha$  implies that the minimum cardinality of a transversal is also  $\alpha$ .)

 $\chi_i := \chi_{\{x_i, y_i\}}$   $(i = 1, ..., \alpha)$ .  $(\chi_i \text{ is 1 in } x_i \text{ and } y_i, \text{ and 0 otherwise.})$  Since every set of H contains exactly one of  $x_i$  and  $y_i$ , the vectors  $\chi_1 - \chi_i$   $(i = 2, ..., \alpha)$  are

G is linearly independent (Padberg), we get that H has at most  $\alpha+1$   $\alpha$ -stable sets.

Comparing this with the trivial lower bound above, we get that  $2\alpha - 2 \le \alpha + 1$ , that is  $\alpha \le 3$ .

Suppose  $\alpha = 3$ . Let S, S', S'' be the three  $\alpha$ -stable sets containing  $u_1$ , where S' is a color class of  $G - u_2$ , and  $S'' = S_1 \cup u_1$ . Let  $T := (S \cup K \cup v_1) \setminus u_1$ .  $|T| = \alpha + \omega - 2$ .

Since by assumption  $S \subseteq S_1 \cup S_2 \cup \{u_1, u_2\}$ , we have  $S \cap S_2 \neq \emptyset$ .  $S \cap S_1 \neq \emptyset$  can also be supposed: if  $S \cap S_1 = \emptyset$ , then  $|S \cap S_2| = 2$ , and  $u_1 S_2 u_2$  is a co-forcing; but then  $u_1 S_1 u_2$  is not a co-forcing (because of (2.2) and the unicity of the  $\alpha$ -stable disjoint from a given  $\omega$ -clique), and interchanging the indices 1 and 2 we have  $S \cap S_1 \neq \emptyset$ ,  $S \cap S_2 \neq \emptyset$ .

It follows that T intersects every  $\alpha$ -stable set except possibly S'. (Note that  $S(K \cup v_1) = S_2 \cup v_2$ .) It also intersects every  $\omega$ -clique except K(S). (The unique clique containing  $u_1$  and not containing  $v_1$  is  $K \cup v_2$ .)

Since  $S' \neq S$ , S' and K(S) are not disjoint. Adding their intersection to T we get a small transversal. This contradiction proves  $\alpha \leq 2$ . Thus G is an antihole.

Let us note that linear algebra can be avoided here at the price of making the proof longer.

### 3. Unique colorability and critical edges

Although the main stream of the paper is to study Problem 2, in this Section we wish to give some particular answers to Problem 1. A general answer can be found in Fonlupt and Sebő (1990), but that cannot be sticked together with Theorem 1.2 to prove SPGC. We would like to provide forcings in uniquely colorable graphs. For some particular graphs we will succeed. These graphs will play a role in the following two sections.

Let us first mention a general relation between the rank r(G) of the set of  $\omega$ -cliques of a graph, perfectness and unique colorability, which is the trivial implication of a characterization of perfectness and unique colorability in Fonlupt, Sebő (1990).

(3.1) Let G be arbitrary. If  $\chi(G) = \omega(G)$ , then  $r(G) \leq n - \omega + 1$ , and if in addition equality holds, then G is uniquely colorable.

Indeed, let  $\chi_1, \ldots, \chi_{\omega}$  be the characteristic vectors of the color classes in an  $\omega$ -coloration. Then the vectors  $\chi_1 - \chi_2, \ldots, \chi_1 - \chi_{\omega}$  are all orthogonal vectors to the  $\omega$ -cliques (like in the proof of Theorem 2.1, Case 2), and they are linearly independent. If there exists another coloration, then let  $\chi_{\omega} + 1$  be a color class in this coloration, which is different

 $\chi_i$   $(i = 1, ..., \omega)$ . It can be easily seen that  $\chi_1 - \chi_\omega + 1$  is linearly independent from  $\chi_2, ..., \chi_1 - \chi_\omega$ , and it is also orthogonal to the  $\omega$ -cliques.

following statement is a key observation for proving theorems 1.3 and 4.1: ) If G is partitionable and K is an  $\omega$ -clique such that G - K is uniquely colorable,

the critical edges of G whose both endpoint is in K form a spanning tree of K. in addition there exists a critical edge as,  $a \in K$ ,  $s \notin K$ , then the critical edges of form a path through all vertices of  $K \cup s$ .

ne of the endpoints of this path is s, let us denote the other by t.  $s(K \setminus t)t$  is a g.

heart of the proof is the following lemma: na: Under the conditions of (3.2) the graph whose vertex set is K and whose edges e critical edges induced by K, is connected.

Let  $\{X_1, X_2\}$  be an arbitrary partition of K, and  $x_1 \in X_1$   $x_2 \in X_2$ . The coloration  $x_1$  consists of S, of  $|X_1| - 1$  color classes intersecting  $X_1$ , and  $|X_2|$  color classes

ecting  $X_2$ . Similarly, the coloration of  $G - x_2$  consists of S, of  $|X_2| - 1$  color classes ecting  $X_2$ , and  $|X_1|$  color classes intersecting  $X_1$ . Thus, since by assumption the tion of these two colorations to G - K is the same, there exists an  $\alpha - 1$ -stable A,  $A \in X_1$ ,  $A \in X_2$  such that  $A \cup A_2$  is a color class of  $A \cup A_1$  is a color of  $A \cup A_2$ . Hence  $A \cup A_2$  is a critical edge.

have proved that there is a critical edge between the classes of any partition of K, proves the lemma.

an be easily proved in various ways that the graph in the lemma cannot contain at. A lemma of Giles, Trotter and Tucker (1984), see also Tucker (1984)) means by that the critical edges form a forest in general. Their proof is tricky but simple, we special case we need is even easier. We use their idea in the following proof:

of (3.2). We first prove (3.2)a: Suppose indirectly that C is a circuit consisting V critical edges,  $V(C) \subseteq K$ . Let  $v_1, \ldots, v_k$  be the vertices of this circuit so that  $v_1, \ldots, v_{n-1}S_{n-1}v_n, v_nS_nv_1$  are co-forcings. It follows from (2.2) that  $S_1 \cup v_2$  is a

lass of  $G - v_1$ . Replace  $S_1 \cup v_2$  by  $S_1 \cup v_1$  in the partition into  $\alpha$ -stable sets ation) of  $G - v_1$ . We get a coloration of  $G - v_2$ . In the same way as before,  $S_2 \cup v_3$  lor class in this coloration.

i+1=1, and finally we get another coloration of  $G-v_1$ . This coloration cannot be the same as the one we started with, because that would mean that we permuted entire color classes. (K cannot contain a color class, because it is a clique.)

Thus the critical edges induced by K form a forest. By the Lemma they form a tree, and a.) is proved.

To prove "b." use (2.1) and the fact that in a tree which is not a path there is no Hamiltonian path: if the statement was not true, s would be joined to any other vertex of K with a path containing at most  $\omega$  vertices; then, according to (2.1), s is adjacent to every vertex of K, and  $K \cup s$  is an  $\omega + 1$ -clique, a contradiction.

Note that the extremities of the path provided by (3.2)b are non-adjacent co-forced vertices. (If they were adjacent,  $K \cup s$  would form an  $\omega + 1$ -clique, see (2.1).)

Let us prove now Theorem 1.3:

Proof of Theorem 1.3 . The conditions of (3.2) are satisfied for both G and its complement  $\bar{G}$ . Apply first (3.2)a. both to G with the clique K, and to  $\bar{G}$  with S.

Let  $s \in S$  be a vertex of degree one in the tree formed by the critical edges of  $\bar{G}$ : the critical co-edges induced by  $S \setminus s$  also form a tree.

Since S is a stable set,  $s \notin Q$  for every forcing of the form  $s_1Qs_2, s_1, s_2 \in S$ . Thus such forcings remain forcings in G-s. Hence the critical co-edges induced by  $S \setminus s$  will be the same in G-s as in G: they form a tree. But forced pairs of vertices have the same color in the perfect graph G-s, implying that  $S\setminus s$  is the subset of a color-class A of G-s. Since A is also an  $\alpha$ -stable,  $A \setminus S$  consists of one vertex, let us denote it by a. as is a critical edge, and by (2.2) K is the unique clique containing a and not containing s: we got in particular that  $a \in K$ .

We can apply (3.2) b.) now: the critical edges induced by  $K \cup s$  form a Hamiltonian path of  $K \cup s$ , where the endpoints of this path are s and t. Thus s and t are co-forced. By (3.2)c.)  $s(K \setminus t)t$  is a forcing, in particular s and t are non-adjacent.

We proved that G has two non-adjacent co-forced vertices. Since the conditions are "self-complementary", we also have two adjacent forced vertices. Thus Theorem 1.1 can be applied.

### 4. $\omega - 1$ -cliques in minimal imperfect graphs

In this section we are going to study  $\omega - 1$ -cliques with the goal of getting closer to forcings. Therefore we are not interested in arbitrary  $\omega - 1$ -cliques, only in those which an participate in forcings, that is, might arise as the intersection of two  $\omega$ -cliques. Unortunately we cannot prove the existence of such  $\omega - 1$ -cliques, that is the existence of forcing (we think this would be a decisive breakthrough concerning Problem 1, and the PGC). Since we want to have existing objects, we define a more general notion, which, the Perfect Graph Conjecture is true is what we want; it defines existing objects; it is estrictive enough to provide an interesting structure.

We suppose all over this section that G is partitionable.

We shall say that an  $\omega - 1$  element subset of an  $\omega$ -clique is an *interval*, if it is a class in partition into  $(\alpha - 1$  pieces of  $\omega - 1$ -cliques of  $G - K \cup S(K)$ , for some  $\omega$ -clique K.

If Q is an  $\omega$ -clique and  $I \subseteq Q$  is an interval, then  $x \in V(G)$  with  $\{x\} = Q \setminus I$  will be led an extremity of I or of Q. The same terminology is used for the intervals of the appearance  $\bar{G}$  of G. (That is an  $\alpha - 1$  element subset of an  $\alpha$ -stable will be called interval, is a color class of a graph of the form  $G - K \cup S(K)$ .)

is easy to check that  $K \setminus x$  is an interval  $(K \text{ is an } \omega\text{-clique}, x \in K)$ , if and only if there its an  $\alpha\text{-stable set } S \text{ with } S \cap K = \{x\}$ , and  $S \cap S(K) \neq \emptyset$ . This will be used without rence.

will turn out that some local structural properties of intervals and of their extremities quivalent to the SPGC.

Every  $\omega - 1$ -clique is contained in at most two  $\omega$ -cliques.

eed, if K is an  $\omega - 1$ -clique, and  $x \in K$  is arbitrary, then in the coloration of G - x are exactly two color classes disjoint from K. Denote these two  $\alpha$ -stables by  $S_1$  and  $\alpha$  other clique than  $K(S_1)$  or  $K(S_2)$  can contain K:

a color class of G - x but  $S \neq S_1$ ,  $S \neq S_2$ , then  $S \cap K \neq \emptyset$ , and K(S) is disjoint this set, whence it does not contain K; if S is not in the coloration of G - x at all,  $\notin K(S)$ , whence it does not contain K again. (4.1) is proved.

that in exactly the same way, every  $\omega - i$ -clique is contained in at most i + 1 es.

Every  $\omega$ -clique has at least two extremities.

Let K be an arbitrary  $\omega$ -clique. Choose  $v \in V(G) \setminus (K \cup S(K))$ . In the coloration

intersecting S(K). These intersect K in different points, which are all extremities of intervals.

The following statement is also an easy exercise, we leave it to the reader.

- (4.3) Every vertex of G is the extremity of at least two  $\omega$ -cliques.
- (4.1) and (4.2) imply that there exist at least n different intervals. (Furthermore, every  $\omega$ -clique can be "represented" by a different interval contained in it ...)

For holes, antiholes (and more generally for "webs") equality holds in these statements. Conversely:

**Theorem 4.1** If G has an  $\omega$ -clique with exactly two extremities, and there is no small transversal in G, then G is a hole or an antihole.

Sketch of the Proof. Let K be an  $\omega$ -clique with exactly two extremities  $a_1, a_2$ . This means exactly, (according to the equivalent definition of intervals noted above), that all  $\alpha$ -stable sets  $T, T \cap S(K) \neq \emptyset, T \neq S(K)$ , contain  $a_1$  or  $a_2$ . In particular, the number of such  $\alpha$ -stable sets is at most  $2(\alpha - 1)$ . (At least one  $\alpha$  stable set containing  $a_1$  is disjoint from S(K): the color class of  $a_1$  in  $G - a_2$ .) Since the  $\alpha$ -stable sets are linearly independent, we get:  $r(\bar{G} - S(K)) \geq n - 2\alpha + 1$ .

Applying (3.1) to  $\bar{G} - S(K)$   $(n - \alpha \text{ replaces } n, \alpha \text{ replaces } \omega)$  we get equality here, and we also get that  $\bar{G} - S(K)$  is uniquely colorable.

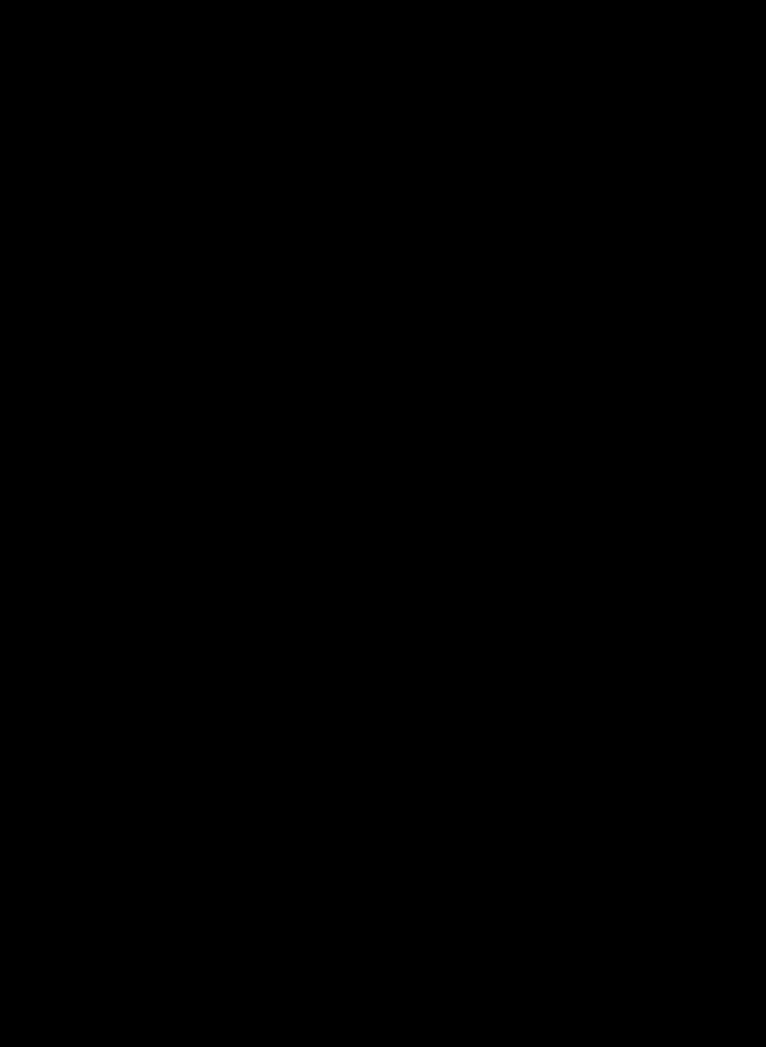
With some work one can also prove that there exists a critical edge from  $a_1$  to S(K). Thus, the conditions of (3.2)b are satisfied, and it follows by (3.2)c that G has adjacent forced vertices.

Now, the condition of Theorem 1.2 is satisfied, and we are done.

Note the similarity between the proof of Theorem 1.3 and 4.1. However, the condition here is not self-complementary, so we have only adjacent forced vertices, and Theorem 1.2 had to be used.

**Theorem 4.2** If G has a vertex which is the extremity of exactly two  $\omega$ -cliques, and there is no small transversal, then G is a hole or an antihole.

Note the kind of polarity relation between (4.2) and (4.3) and between Theorem 4.1 and 4.2. (This can be given a precise sense: Tucker has noted that the intersection graph of a partitionable graph is also partitionable.) However, we cannot prove one statement from the other, and Theorem 4.1 seems to be much more difficult: we use Theorem 1.2 in an essential way in the proof, whereas this is not necessary in the (here omitted) proof of



Theorem 4.2. The reason why we cannot prove one of these statements from the other is that we do not know whether the polarity relation keeps small transversals.

### 5. Reformulations of the Strong Perfect Graph Conjecture

The goal of this section is to collect new reformulations of the SPGC which follow from the results of the paper.

We distinguish two levels of reformulations: to prove the SPGC from the second level statements, Theorem 1.2 seems to be essential, whereas the first level statements imply SPGC more easily.

Anyone of the following statements is equivalent to the Strong Perfect Graph Conjecture:

### 1st level

- If G is minimal imperfect, then for every  $\omega$ -clique K, G-K is uniquely colorable;
- If G is minimal imperfect, then the number of intervals is |V(G)|;
- f G is minimal imperfect, then the set of intervals is linearly independent;
- f G is minimal imperfect, then every interval has two extremities;
- f G is minimal imperfect, then every  $\omega$ -clique has two extremities;
- fG is minimal imperfect, then every vertex is the extremity of two intervals;
- G is minimal imperfect, then there exists a vertex which is the extremity of two intervals;

### nd level

- G is minimal imperfect, then there exists an  $\omega$ -clique K, such that  $G (K \cup S(K))$  and s complement are uniquely colorable;
- G is minimal imperfect, then there exist two adjacent forced vertices (or more generally, vo co-forcings and one forcing like in Theorem 2.1) in G or  $\bar{G}$ ;
- G is minimal imperfect, then there exists a forcing xKy such that  $G (K \cup \{x,y\})$  is niquely colorable;
- G is minimal imperfect, then there exists a forcing xKy such that every clique intersectg K also intersects x or y.
- G is minimal imperfect, then there exists an interval which has exactly two extremities.
- Can any of these be proved for a subclass for which the SPGC is open?
- cknowledgment: I am indebted to Frédéric Maffray and Myriam Preissmann for many ecious comments.

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