

## 4. Complements to the first 3 series

# The postman polyhedron

**Def :**  $\delta(W) \subseteq E(G)$  ( $W \subseteq V$ ) is a *T-cut*, if  $|W \cap T|$  is odd

**Proposition :**  $F$  T-join,  $\delta(W)$  T-cut  $\Rightarrow |F \cap \delta(W)| \geq 1$

**Theorem** Edmonds, Johnson (1973) :  $Q_+(G, T) := \text{conv}(\text{T-joins}) + \mathbb{R}_+^n =$   
 $\{x \in \mathbb{R}_+^E \mid x(\delta(W)) \geq 1, \delta(W) \text{ is a T-cut, i.e. } |W \cap T| \text{ is odd}\}$

**Proof :** Through the following slides.

# Minmax bipartite

$$\tau(G,T) \quad := \quad \min \{ |F| : F \subseteq E, F \text{ is a T-join} \}$$

$$\nu(G,T) \quad := \quad \max \{ |\mathcal{C}| : \mathcal{C} \text{ disjoint T-cuts} \}$$

Easy :  $\tau(G,T) \geq \nu(G,T)$

**Theorem** (Seymour '81) If  $G$  is bipartite,

$$\tau(G,T) = \nu(G,T)$$

# Minmax nonbipartite

$v_2(G,T) := \max\{ |\mathcal{C}| : \mathcal{C} \text{ 2-packing of T-cuts} \}$ , where a *2-packing* is a family covering every element  $\leq$  twice

**Easy** :  $\tau(G,T) \geq v_2(G,T) / 2$

**Proof** : Let  $F$  be a T-join, and  $\mathcal{C}$  a 2-packing of T-cuts.

Then  $2 |F| = \sum_{C \in \mathcal{C}} |F \cap C| \geq v_2(G,T)$

**Theorem** (Edmonds-Johnson '73) If  $G$  is arbitrary,

$$\tau(G,T) = v_2(G,T) / 2$$

# Packing

A *packing* is a family covering every element  $\leq$  once

A *2-packing* is a family covering every element  $\leq$  twice

$$v_2(G, T) / 2 \geq v(G, T)$$

(Possibly) fractional : coefficients  $y_C$  ( $C \in \mathcal{C}$ ) whose sum has to be maximized :  $v^*$  for packings .

For  $c: E \rightarrow \mathbb{R}_+$  :  $v(G, T, c)$  ,  $v_2(G, T, c)$  ,  $v^*(G, T, c)$

# Linear Programming

## Duality Theorem

$$Ax \leq b$$

$$(A \in \mathbb{Q}^{m \times n}, b, c \in \mathbb{Q}^n)$$

$$\max c^T x$$

dual:

=

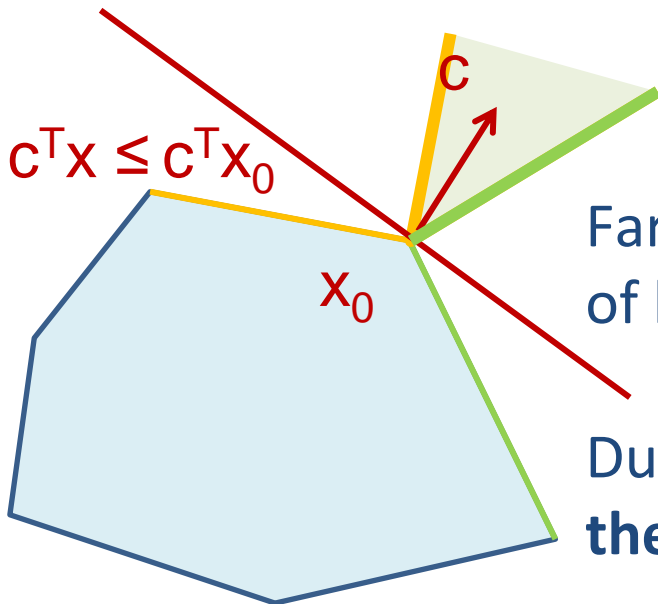
$$yA = c$$

$$y \geq 0$$

$$\min y^T b$$

Weak duality

$\leq$  : every primal  $x$  is feasible for all nonnegative combinations



Farkas Lemma: Every 'tight' consequence of linear inequalities is their nonneg.lin.comb.

Duality Theorem  $\Leftrightarrow$  If the max for  $c$  is  $c^T x_0$ , then  $cx \leq c^T x_0$  is a nonneg.lin.comb of  $a_i^T x_0 \leq b_i$

# Weak duality for the T-join polyhedron

Let  $F$  be a T-join, and  $\mathcal{C}$  a 2-packing of T-cuts.

$$\text{Then } 2 |F| = \sum_{C \in \mathcal{C}} |F \cap C| \geq v_2(G, T)$$

Let  $F$  be a T-join, and  $\mathcal{C}$  a (possibly fractional) 1-packing of T-cuts with coefficients  $y_C$  ( $C \in \mathcal{C}$ )

$$\text{Then } |F| = \sum_{C \in \mathcal{C}} y_C |F \cap C| \geq v^*(G, T)$$

Let  $F$  be a T-join, and  $\mathcal{C}$  a (possibly fractional)  $c$ -packing of T-cuts with coefficients  $y_C$  ( $C \in \mathcal{C}$ )

$$\text{Then } |F| = \sum_{C \in \mathcal{C}} y_C |F \cap C| \geq v^*(G, T, c) \text{ (or } v(G, T, c))$$

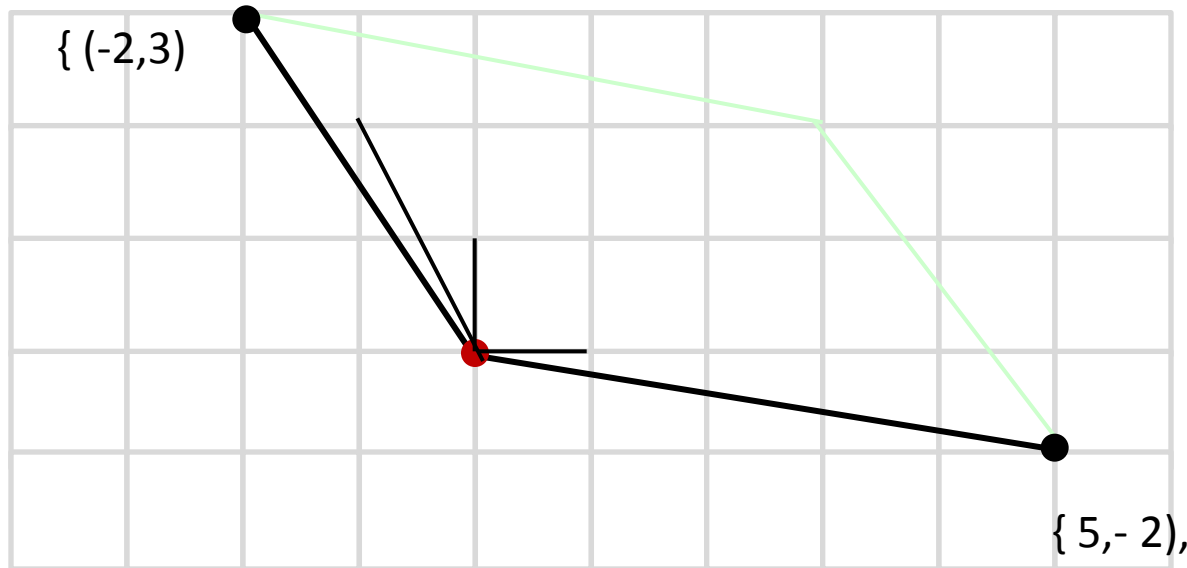
# Linear Programming

Hilbert bases (normal semigroups)

$H \subseteq \mathbb{Z}^n$  is a *Hilbert basis* if any integer vector which is a **nonneg** comb is also a **nonneg integer** comb

Example  $n=2$  :

pointed cone : cone  $\{ (-2,3), (5,-1) \}$   $|\det| = 13$



adding  $\{ (-1,2), (0,1), (1,0) \}$  : Hilbert basis

*Integer Caratheodory property* (+‘partition’ into unimodular cones)



# Proving the T-join polyhedron Thm

$$Q_+(G,T) = \{x \in \mathbb{R}^E : x(W) \geq 1, W \text{ is a T-cut}, x \geq 0\}$$

Edmonds-Johnson:  $\frac{1}{2}$  TDI, vertices: T-joins



$$\tau(G,T, c) = v(G,T, 2c) / 2$$



Seymour (81): If  $G$  arbitrary,  $\tau(G,T) = v_2(G,T) / 2$



Edmonds, Johnson(73): If  $G$  is bipartite,  $\tau(G,T) = v(G,T)$

**Metatheorem** : Polyhedron the same as weighted minmax theorem

# If negative weights are allowed ?

$$c(F) = |c|(F \setminus E_-) - |c|(F \cap E_-) = |c|(F \Delta E_-) - c(E_-)$$

(So if  $(G,w)$  is conservative,  $\lambda_w(x,y) := \min \{w(P) : P \text{ path}\} =$   
 $\min \{w(P) : P \{x,y\}\text{-join}\}$

Is reducible to min weight perfect matchings.)

This reduction leads to the  
T-join *polytope*

# Another application

## SCHEDULING IDENTICAL JOBS ON 2 IDENTICAL MACHINES

**Input:** Partially ordered set of tasks of unit length.

**Output:** Schedule of min completion time  $T$

**Theorem** : (Fujii & als) :  $T = n - \nu(G_{\text{input}})$

Solutions for max (weighted) matchings:

with Edmonds' algorithm (1965)

Grötschel, Lovász, Schrijver

with Padberg-Rao (1979)



To come : matroids

**Exercises to revise for the third course : series 7.**

